

# MEMORY SYSTEM

Mahdi Nazm Bojnordi

Assistant Professor

School of Computing

University of Utah

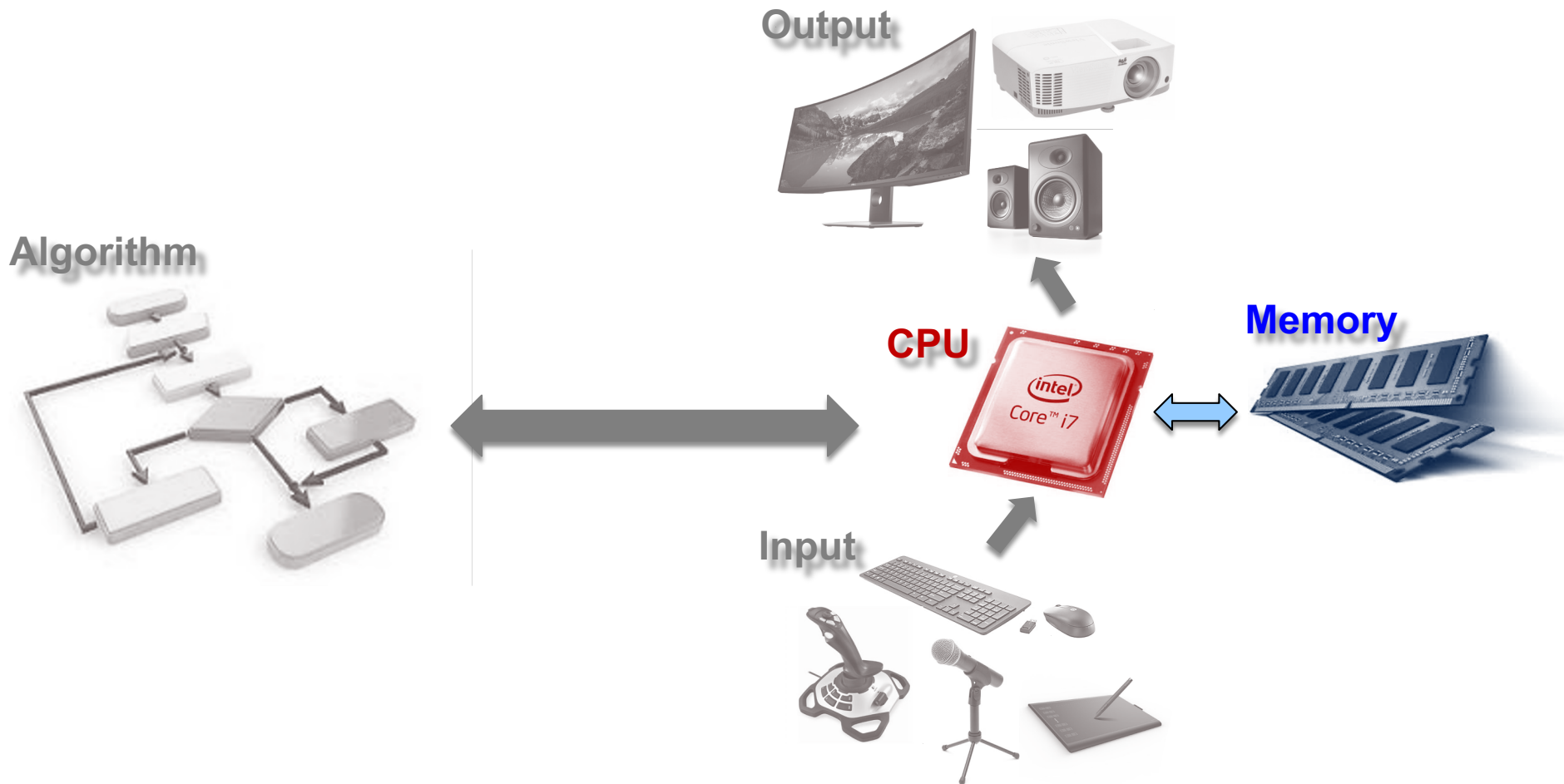
# Overview

---

- This lecture
  - ▣ Memory system
    - Cache

# Computer Organization

- Classic components of a computing system



# Memory System

- Data and instructions are stored on DRAM chips
  - ▣ DRAM has high bit density and low speed
  - ▣ An access DRAM may take about 300 processor cycles

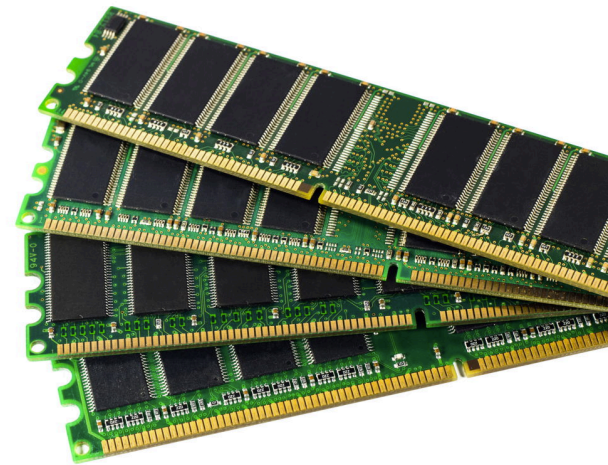
**Processor**



~300X

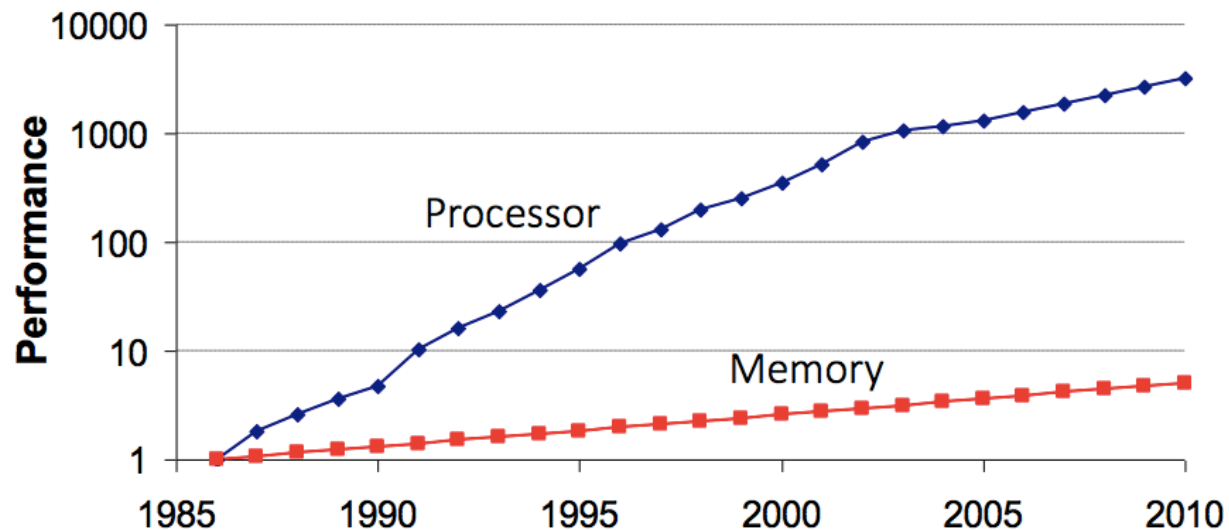


**Memory**



# The Memory Wall

- Processor-memory performance gap increased over 50% per year
  - ▣ Processor performance historically improved  $\sim 60\%$  per year
  - ▣ Main memory access time improves  $\sim 5\%$  per year



# Memory System

- Data and instructions are stored on DRAM chips
  - ▣ DRAM has high bit density and low speed
  - ▣ An access DRAM may take about 300 processor cycles
- **How to bridge the speed gap?**

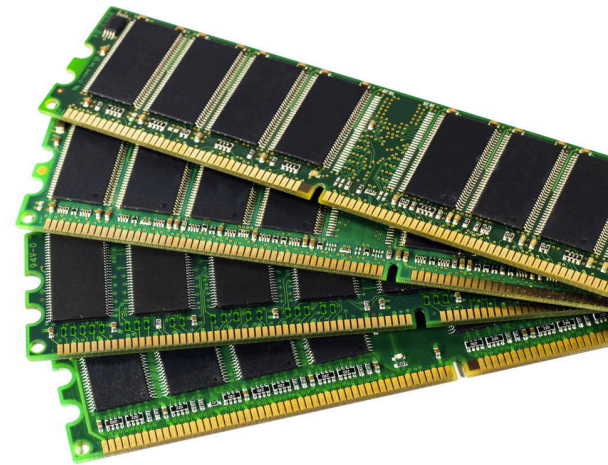
Processor



~300X

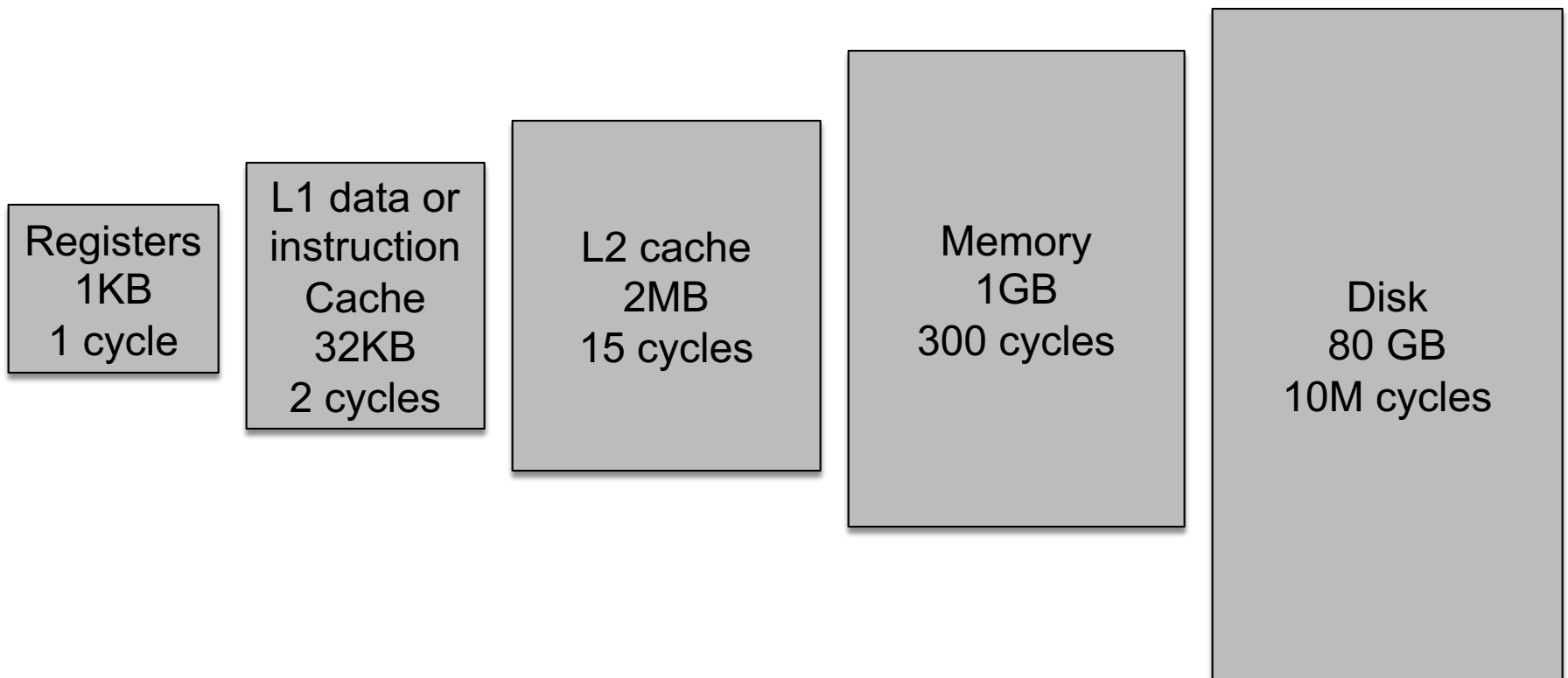


Memory



# Memory Hierarchy

- The basic structure of a memory hierarchy.



# Processor Cache

- Occupies a large fraction of die area in modern microprocessors

3-3.5 GHz  
~\$1000 (2014)





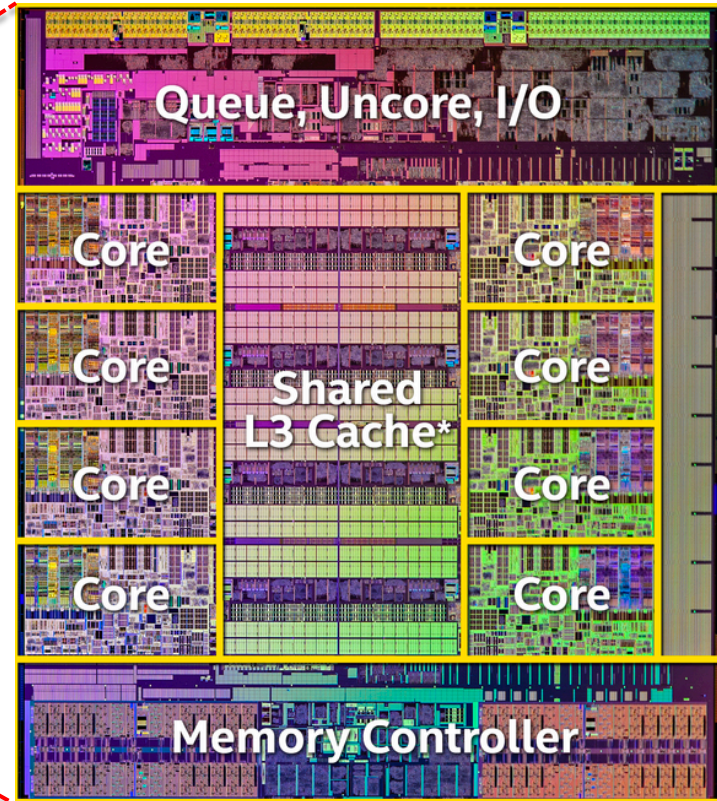
# Processor Cache

- Occupies a large fraction of die area in modern microprocessors

3-3.5 GHz  
~\$1000 (2014)



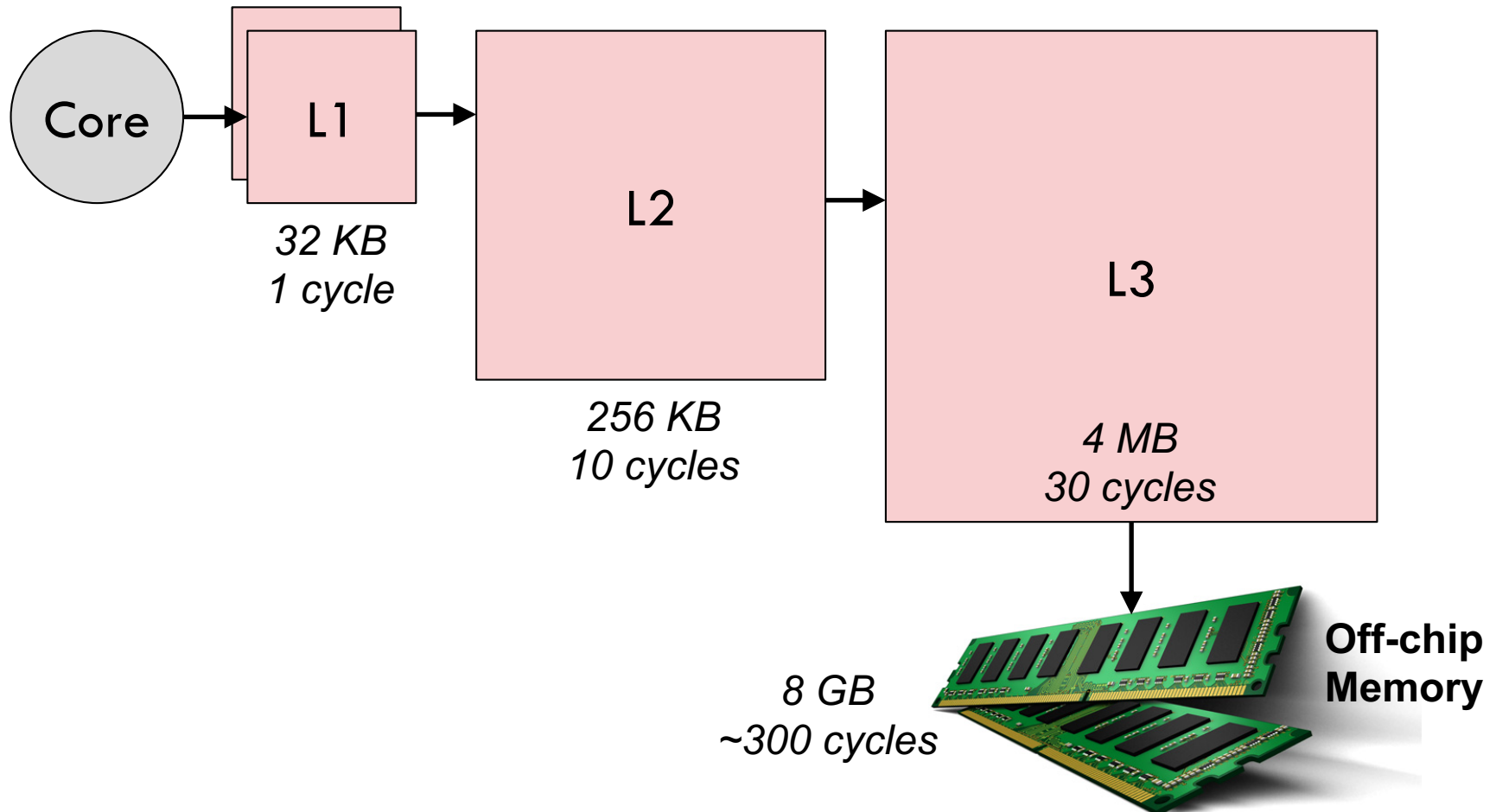
20MB of cache



Source: Intel Core i7

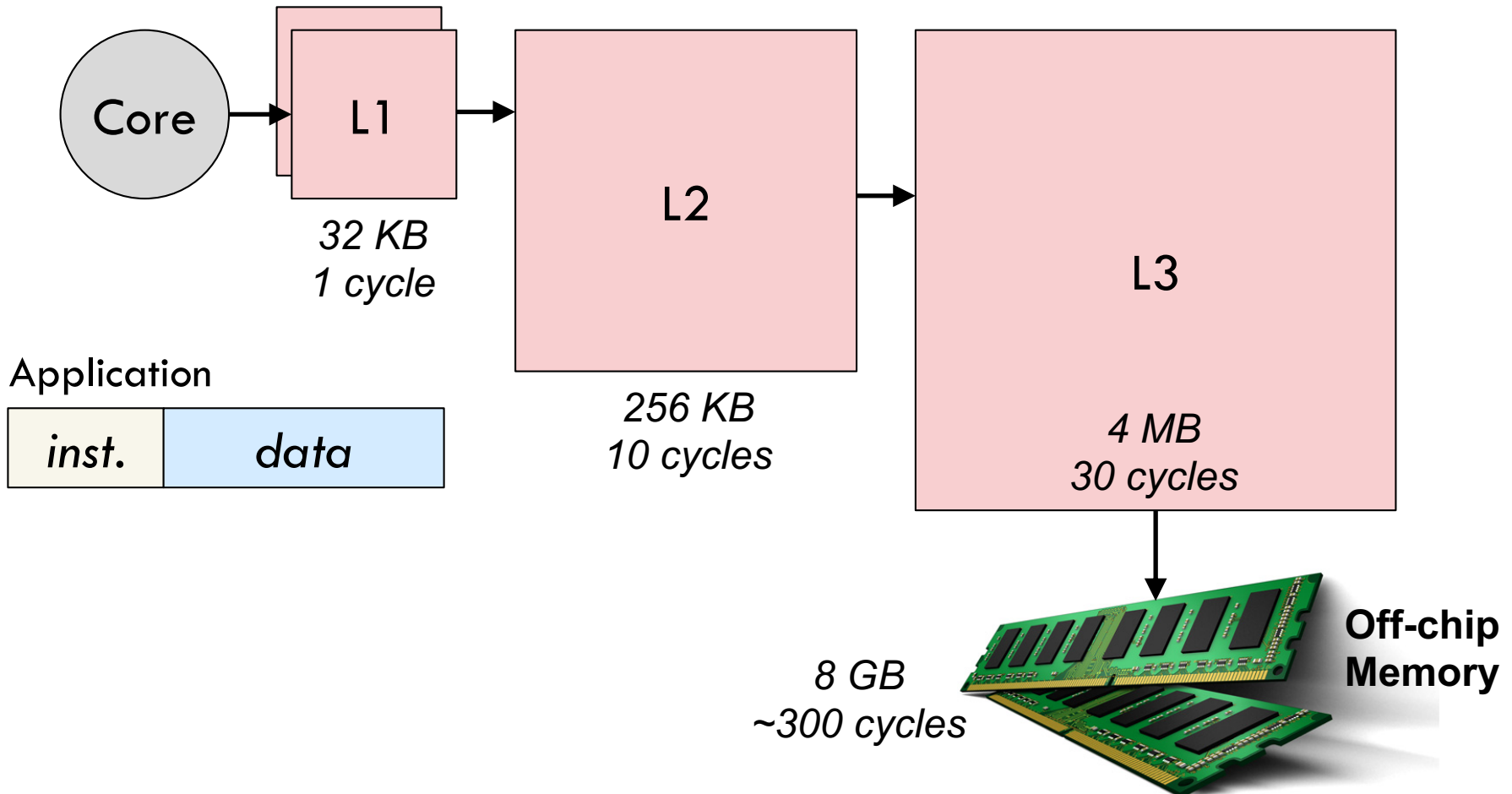
# Cache Hierarchy

- Example three-level cache organization



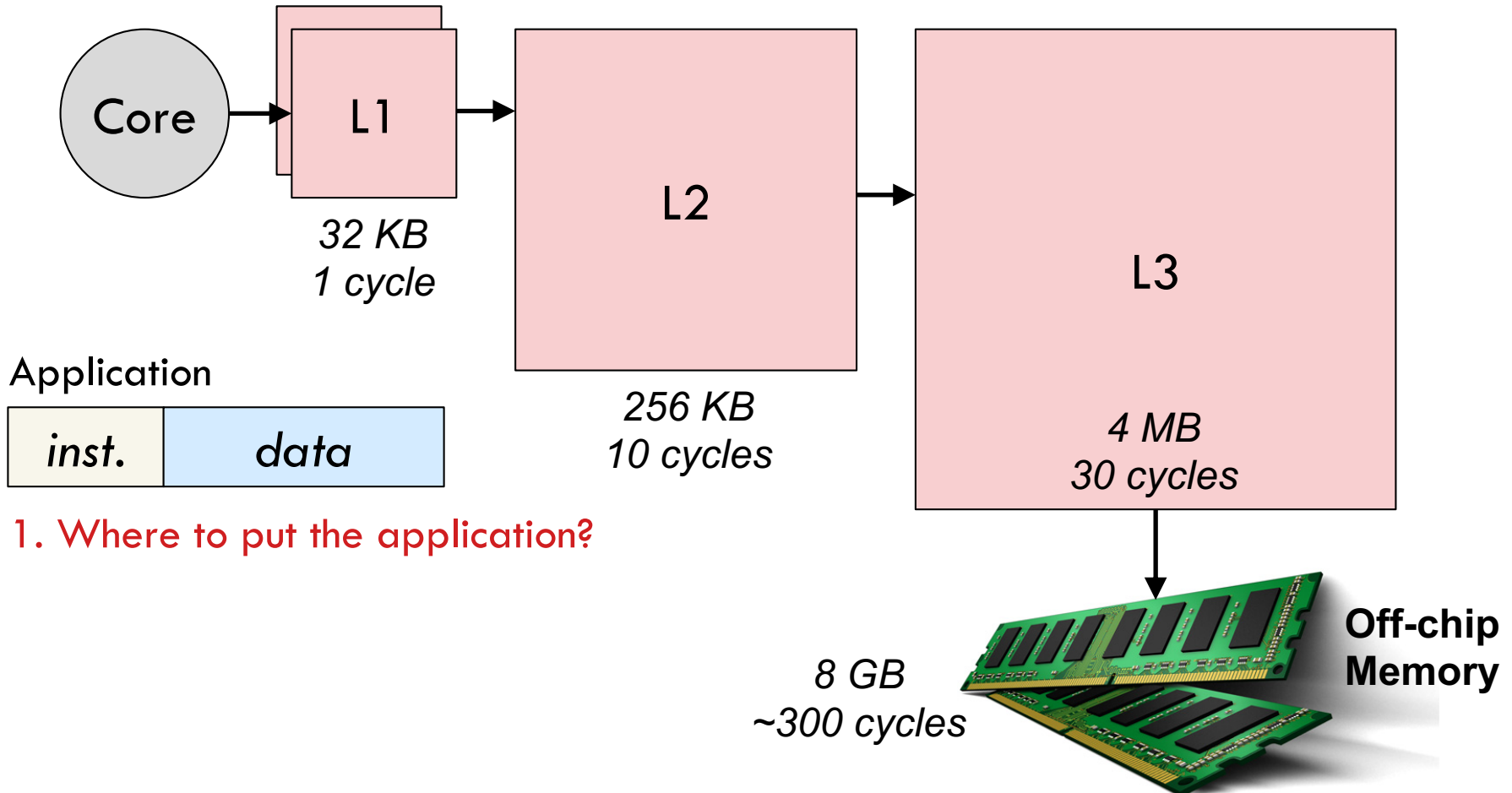
# Cache Hierarchy

- Example three-level cache organization



# Cache Hierarchy

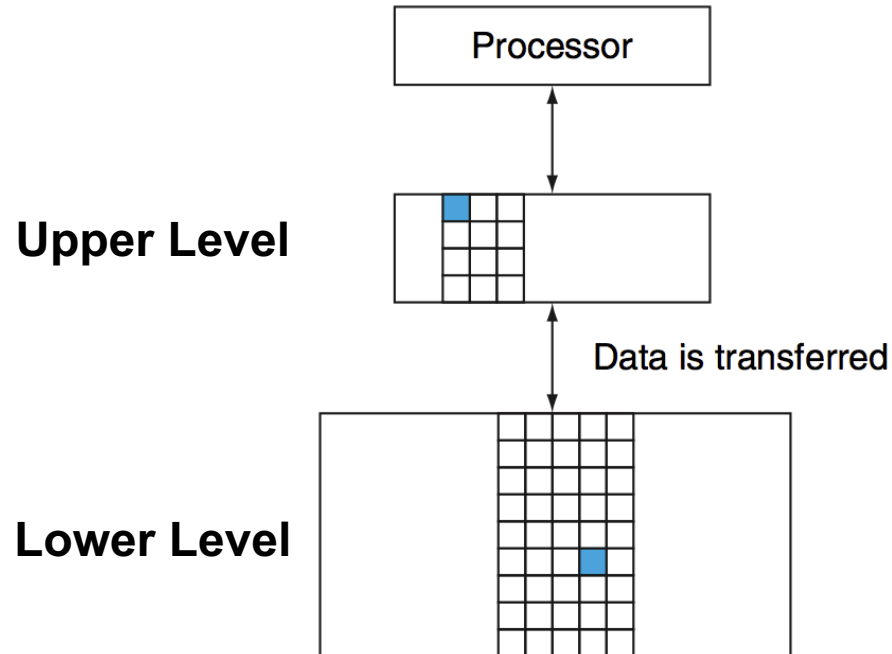
- Example three-level cache organization



# Memory Hierarchy

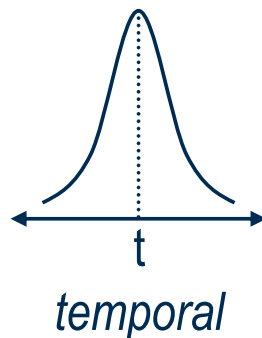
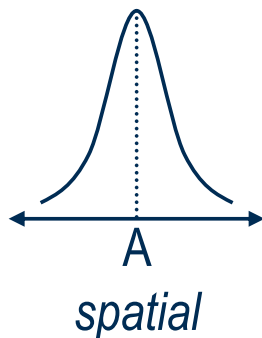
- The basic structure of a memory hierarchy.
- Multiple levels of the memory

**Idea: keep important data closer to processor.**



# Principle of Locality

- Memory references exhibit localized accesses
- Types of locality
  - ▣ *spatial*: probability of access to  $A+\delta$  at time  $t+\varepsilon$  highest when  $\delta \rightarrow 0$
  - ▣ *temporal*: probability of accessing  $A+\varepsilon$  at time  $t+\delta$  highest when  $\delta \rightarrow 0$

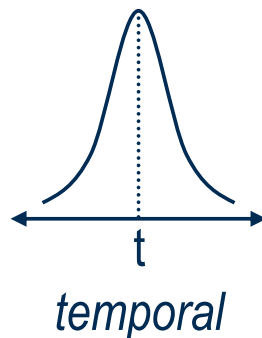
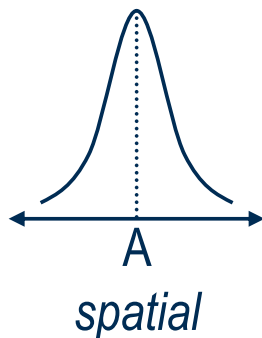


```
for (i=0; i<1000; ++i) {  
    sum = sum + a[i];  
}
```

Key idea: store local data in fast cache levels

# Principle of Locality

- Memory references exhibit localized accesses
- Types of locality
  - ▣ *spatial*: probability of access to  $A+\delta$  at time  $t+\varepsilon$  highest when  $\delta \rightarrow 0$
  - ▣ *temporal*: probability of accessing  $A+\varepsilon$  at time  $t+\delta$  highest when  $\delta \rightarrow 0$



```
for (i=0; i<1000; ++i) {  
    sum = sum + a[i];  
}
```

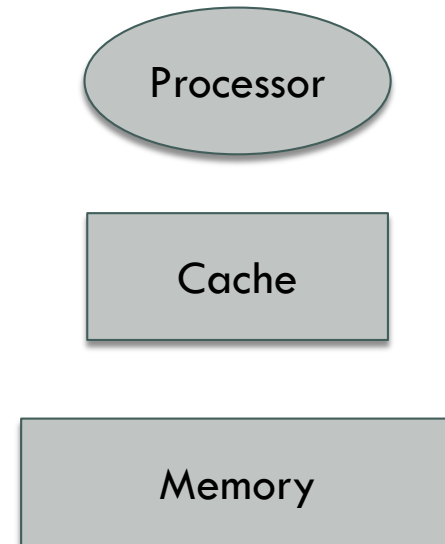
**temporal**      **spatial**

Diagram illustrating the code snippet with annotations. The code is: `for (i=0; i<1000; ++i) { sum = sum + a[i]; }`. A red arrow points from the word **temporal** to the `sum` variable in the assignment statement. Another red arrow points from the word **spatial** to the `a[i]` array access in the same statement.

Key idea: store local data in fast cache levels

# Cache Architecture

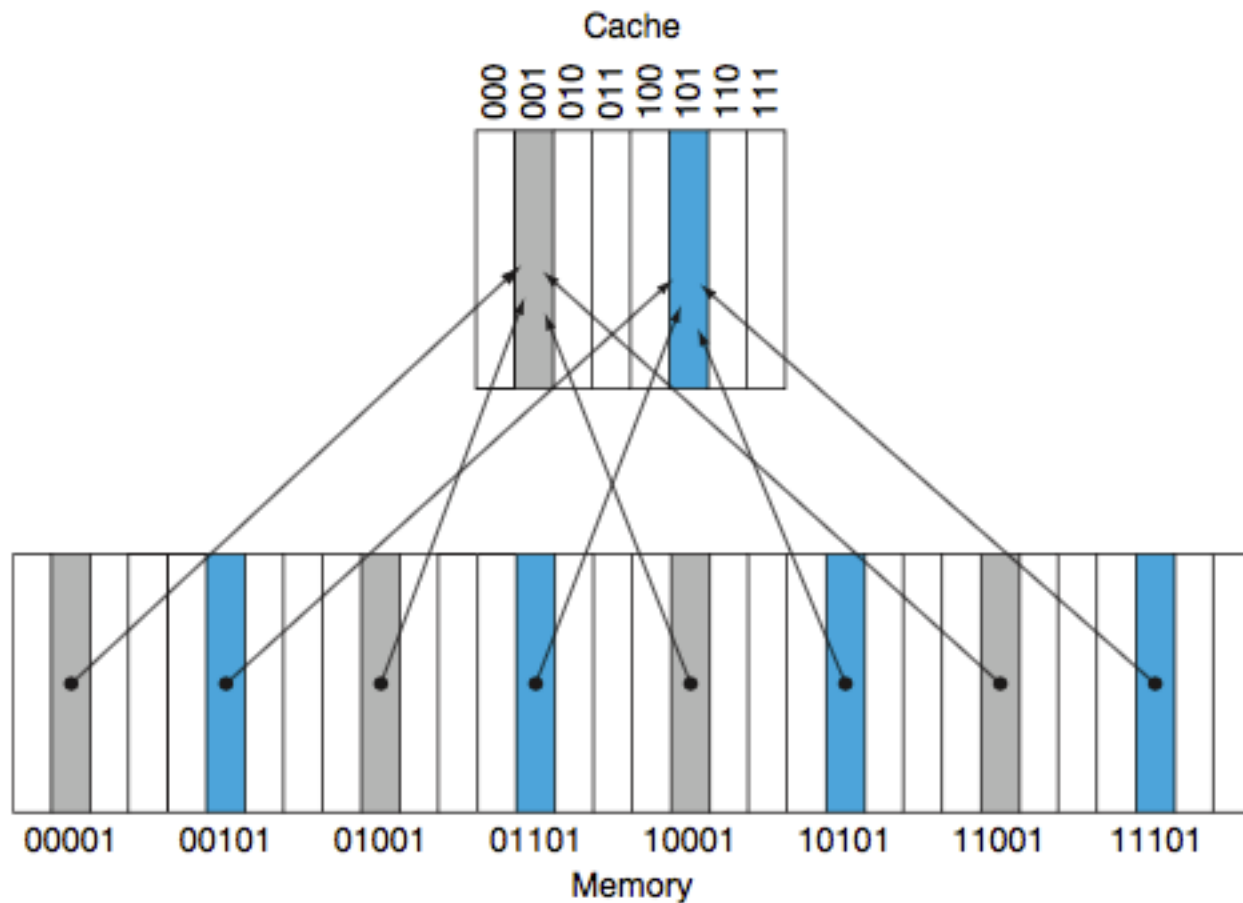
- Design principles
  - ▣ Temporal locality: if you used some data recently, you will likely use it again
  - ▣ Spatial locality: if you used some data recently, you will likely access its neighbors
- Cache terminology
  - ▣ Access time
  - ▣ Hit vs. miss
  - ▣ Miss penalty





# Direct-Mapped Cache

- Cache address



# Direct-Mapped Cache

- Cache lookup

