

HARDWARE FOR ARITHMETIC

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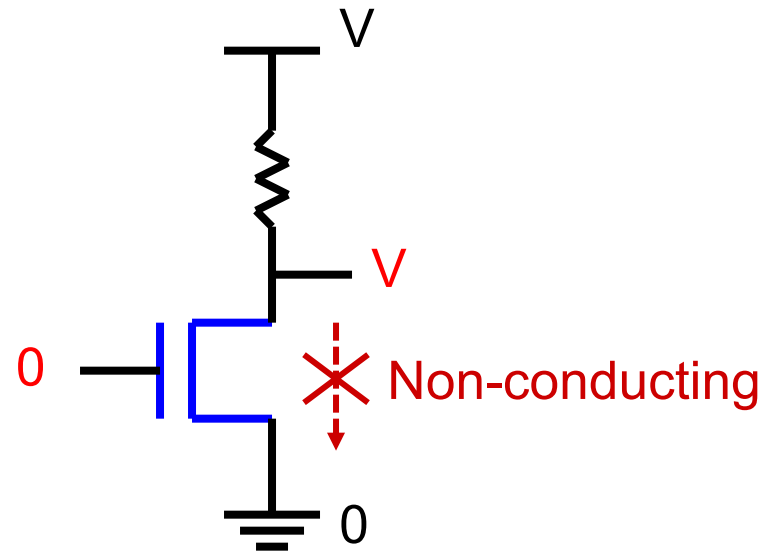
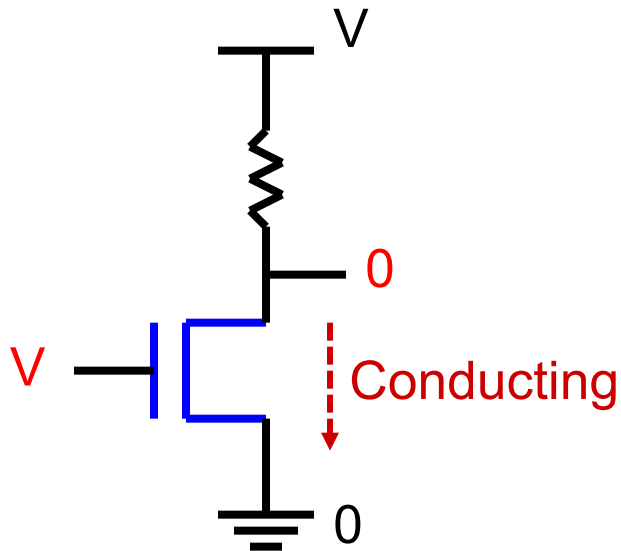
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Overview

- Past lectures
 - ▣ High-level stuff mostly on ISA, Assembly, and number representation
- This lecture
 - ▣ Basics of logic design
 - ▣ Hardware for arithmetic

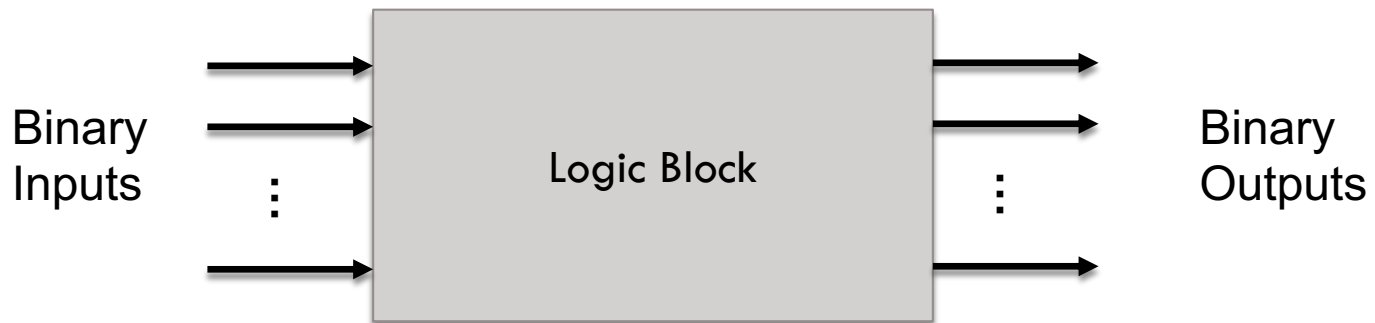
Fundamentals of Digital Design

- **Binary logic:** two voltage levels
 - ▣ high and low; 1 and 0; true and false
- **Binary arithmetic**
 - ▣ Based on a 3-terminal device that acts as a switch



Logic Blocks

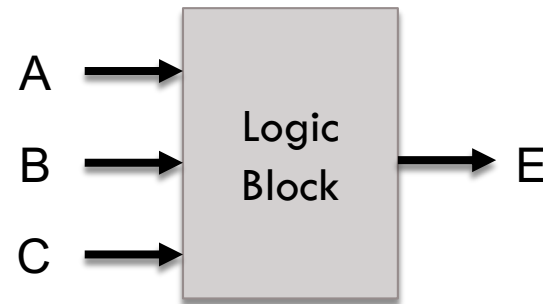
- A logic block comprises binary inputs and binary outputs
 - ▣ **Combinational**: the output is only a function of the inputs
 - ▣ **Sequential**: the block has some internal memory (state) that also influences the output
- **Gate**: a basic logic block that implements AND, OR, NOT, etc.



Logic Blocks: Truth Table

- A **truth table** defines the outputs of a logic block for each set of inputs
 - ▣ **Example:** consider a block with 3 inputs A, B, C and an output E that is true only if exactly 2 inputs are true

A	B	C	E
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	0



Boolean Algebra

- Three primary operators are used to realize Boolean functions

- Boolean operations
 - ▣ OR (symbol +)
 - $X = A + B$: X is true if at least one of A or B is true
 - ▣ AND (symbol .)
 - $X = A . B$: X is true if both A and B are true
 - ▣ NOT (symbol $\bar{\quad}$)
 - $X = \bar{A}$: X is the inverted value of A

Pictorial Representation

- Logic gates

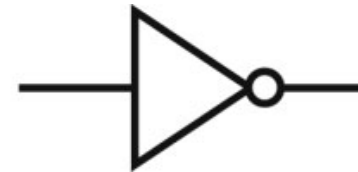
AND



OR



NOT



- What function is the following?



Boolean Algebra Rules

- Identity law

- ▣ $A + 0 = A$

- ▣ $A \cdot 1 = A$

- Zero and One laws

- ▣ $A + 1 = 1$

- ▣ $A \cdot 0 = 0$

- Inverse laws

- ▣ $A \cdot \overline{A} = 0$

- ▣ $A + \overline{A} = 1$

- Commutative laws

- ▣ $A + B = B + A$

- ▣ $A \cdot B = B \cdot A$

- Associative laws

- ▣ $A + (B + C) = (A + B) + C$

- ▣ $A \cdot (B \cdot C) = (A \cdot B) \cdot C$

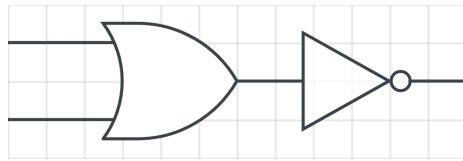
- Distributive laws

- ▣ $A \cdot (B + C) = (A \cdot B) + (A \cdot C)$

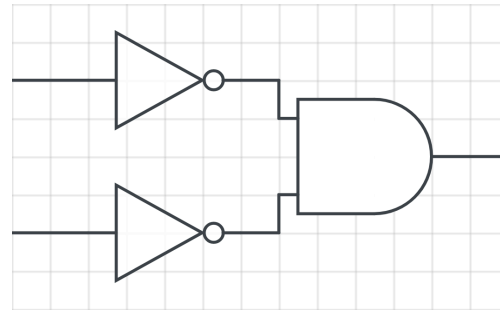
- ▣ $A + (B \cdot C) = (A + B) \cdot (A + C)$

DeMorgan's Law

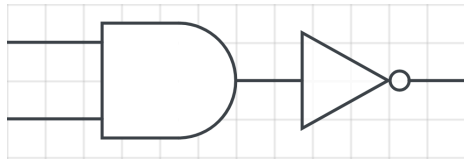
□ $\overline{A + B} = \overline{A} \cdot \overline{B}$



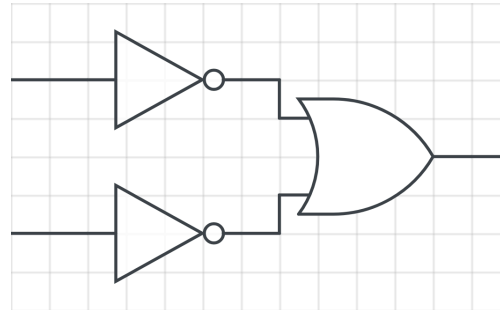
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□ $\overline{A \cdot B} = \overline{A} + \overline{B}$



=



Example: Boolean Equation

- Consider the logic block that has an output E that is true only if exactly two of the three inputs A, B, C are true

- Multiple correct equations
 - Two must be true, but all three cannot be true
 - $E = ((A \cdot B) + (B \cdot C) + (A \cdot C)) \cdot \overline{(A \cdot B \cdot C)}$

 - Identify the three cases where it is true
 - $E = (A \cdot B \cdot \overline{C}) + (A \cdot C \cdot \overline{B}) + (C \cdot B \cdot \overline{A})$

Implementing Boolean Functions

- Can realize any logic block with the AND, OR, NOT
 - ▣ Draw the truth table
 - ▣ For each true output, represent the corresponding inputs as a product
 - ▣ The final equation is a sum of these products

A	B	C	E
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	0

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$\overline{A}.B.C$

$A.\overline{B}.C$

$A.B.\overline{C}$

Implementing Boolean Functions

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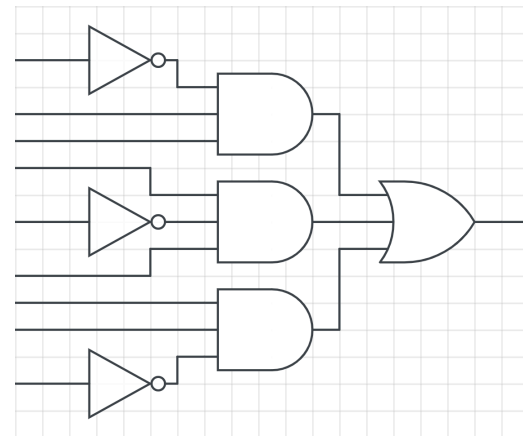
$\bar{A}.B.C$

$A.\bar{B}.C$

$A.B.\bar{C}$

Sum of Products

$$E = (\bar{A}.B.C) + (A.\bar{B}.C) + (A.B.\bar{C})$$

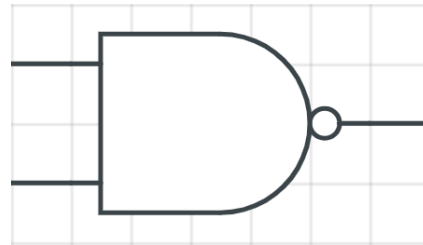


Universal Gates

- Universal gate is a logic that can be used to implement any complex function

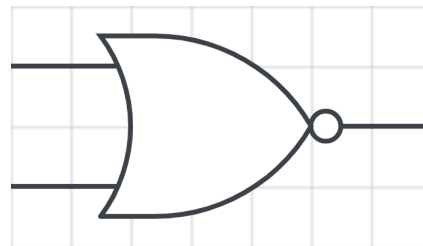
- NAND

- Not of AND
- $A \text{ nand } B = \overline{A \cdot B}$



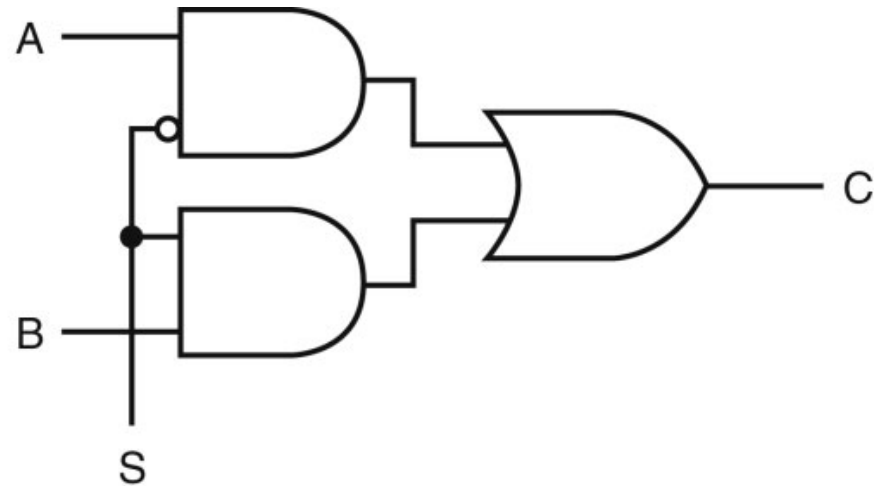
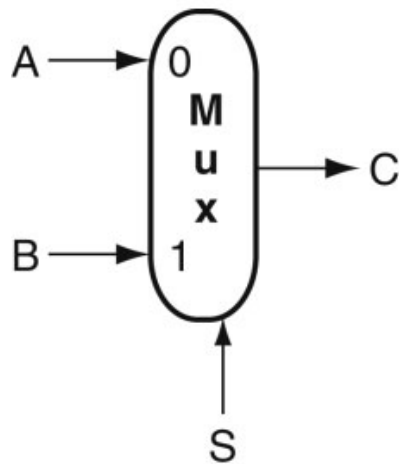
- NOR

- Not of OR
- $A \text{ nor } B = \overline{A + B}$



Common Logic Block

- A **multiplexer (or selector)** reflects one of n inputs on the output depending on the value of the select bits
 - ▣ Example: 2-input mux



Common Logic Block

- A **full adder** generates the sum and carry for each bit position

A	B	C _{in}	Sum	C _{out}
0	0	0		
0	0	1		
0	1	0		
0	1	1		
1	0	0		
1	0	1		
1	1	0		
1	1	1		

$$\begin{array}{r} 1 \quad 0 \quad 0 \quad 1 \\ 0 \quad 1 \quad 0 \quad 1 \\ \hline \end{array}$$

Sum
Cout

Common Logic Block

- A **full adder** generates the sum and carry for each bit position

A	B	C _{in}	Sum	C _{out}
0	0	0		
0	0	1		
0	1	0		
0	1	1		
1	0	0		
1	0	1		
1	1	0		
1	1	1		

	1	0	0	1
	0	1	0	1
	<hr/>			
Sum	1	1	1	0
Cout	0	0	0	1

Common Logic Block

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A	B	C _{in}	Sum	C _{out}
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

	1	0	0	1
	0	1	0	1
Sum	1	1	1	0
Cout	0	0	0	1

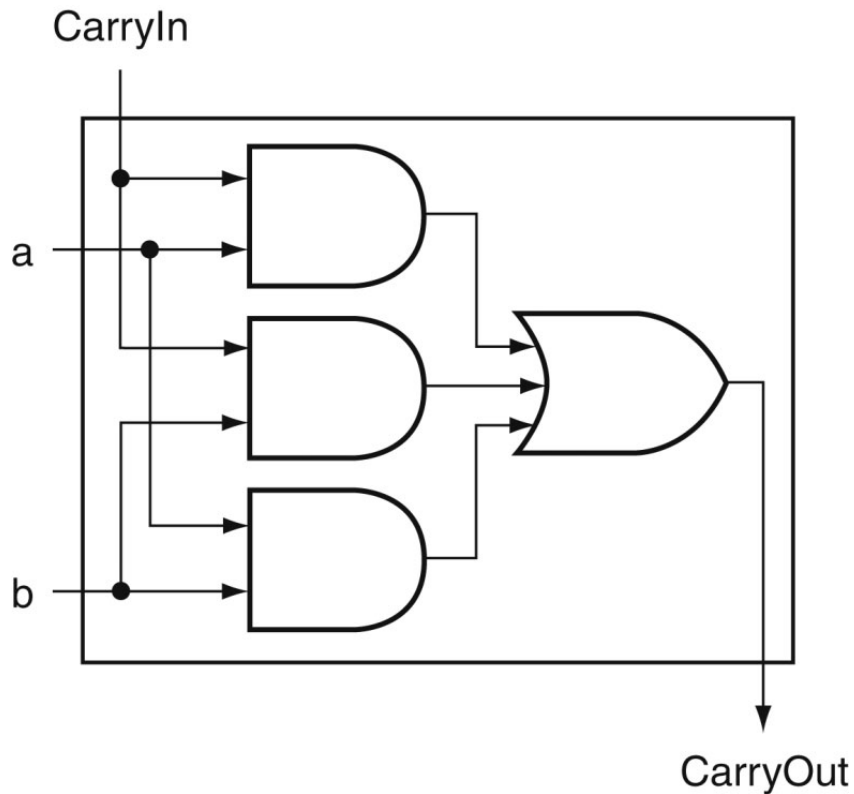
Equations:

$$\text{Sum} = C_{in} \cdot \bar{A} \cdot \bar{B} + B \cdot \bar{C}_{in} \cdot \bar{A} + A \cdot \bar{C}_{in} \cdot \bar{B} + A \cdot B \cdot C_{in}$$

$$\text{Cout} = A \cdot B \cdot C_{in} + A \cdot B \cdot \bar{C}_{in} + A \cdot C_{in} \cdot \bar{B} + B \cdot C_{in} \cdot \bar{A} = A \cdot B + A \cdot C_{in} + B \cdot C_{in}$$

Common Logic Block

- A **full adder** generates the sum and carry for each bit position



	1	0	0	1
	0	1	0	1
<hr/>				
Sum	1	1	1	0
Cout	0	0	0	1

Equations:

$$\text{Sum} = C_{in} \cdot \bar{A} \cdot \bar{B} + B \cdot \bar{C}_{in} \cdot \bar{A} + A \cdot \bar{C}_{in} \cdot \bar{B} + A \cdot B \cdot C_{in}$$

$$\begin{aligned} \text{Cout} = & A \cdot B \cdot C_{in} + A \cdot B \cdot \bar{C}_{in} + A \cdot C_{in} \cdot \bar{B} + B \cdot C_{in} \cdot \bar{A} = \\ & A \cdot B + A \cdot C_{in} + B \cdot C_{in} \end{aligned}$$